

Symmetric I* Restriction Method of Fuzzy Inference

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Abstract: As one of important parts of fuzzy logic, fuzzy inference plays a vital role in the fields of fuzzy control, artificial intelligence, affective computing, image processing and so forth. Two key problems of fuzzy inference are FMP (fuzzy modus ponens) and FMT (fuzzy modus tollens). How to get the ideal solution for FMP and FMT is a difficult problem in the area of fuzzy logic. Aiming at such problem, from the idea of symmetric implicational reasoning, triple I* method and restriction theory, we put forward and investigate the α -symmetric I* restriction method, and then generalize it to the $\alpha(x,y)$ -symmetric I* restriction method. To begin with, the α -symmetric I* restriction principle and the $\alpha(x,y)$ -symmetric I* restriction principle are established. Furthermore, the equivalent condition to let a basic restriction solution exist is given. Then the unified solutions of the α -symmetric I* restriction method and the $\alpha(x,y)$ -symmetric I* restriction method are achieved for R-implications and (S, N)-implications. Besides, some special cases of optimal solutions are shown. Finally, the corresponding conclusions are provided when the two methods degenerate into the α -triple I* restriction method and $\alpha(x,y)$ -triple I* restriction method. These research results would be an important improvement for the fields of fuzzy inference, fuzzy logic and related applications.

Keywords: Fuzzy Inference, Fuzzy Implication, Triple I Method, Symmetric Implicational Method

1. Introduction

Nowadays fuzzy inference plays a vital role in the fields of fuzzy control, artificial intelligence, affective computing, image processing and so on [1-5]. Two key problems of fuzzy inference are FMP (fuzzy modus ponens) and FMT (fuzzy modus tollens) [6-7] denoted as follows:

FMP: for $A \rightarrow B$ and A^* , to compute B^* , (1)

FMT: for $A \rightarrow B$ and B^* , to compute A^* . (2)

Here $x \in X$ (the input universe), $y \in Y$ (the output universe), and $A, A^* \in F(X)$ (the set of all fuzzy subsets on X), $B, B^* \in F(Y)$ (the set of all fuzzy subsets on Y). For this field, the most classical algorithm is the CRI (compositional rule of inference) method [8-10].

To get better results, Wang [11] proposed the triple I method.

Its ideal solution was the smallest $B^* \in F(Y)$ (or the largest $A^* \in F(X)$) such that

$$(A(x) \rightarrow B(y)) \rightarrow (A^*(x) \rightarrow B^*(y)) \quad (3)$$

is maximized for any $x \in X, y \in Y$, where \rightarrow employs a fuzzy implication. Following that, Song et al. [12-13] established the triple I restriction method, whose ideal solution was the largest $B^* \in F(Y)$ (or the smallest $A^* \in F(X)$) such that ($\alpha \in (0, 1]$)

$$(A(x) \rightarrow B(y)) \rightarrow (A^*(x) \rightarrow B^*(y)) < \alpha \quad (4)$$

holds for any $x \in X, y \in Y$.

Later, a lot of scholars carried through researches related to the triple I method and the triple I restriction method [14-15]. Tang et al. proposed the universal triple I method [16]. From the viewpoints of both fuzzy system and fuzzy reasoning, the

α -universal triple I restriction method was proposed and investigated [17]. The variable differently implicational algorithm was put forward, which made the current differently implicational algorithms compose a united whole [18]. The main condition of the differently implicational inference algorithm was reconsidered from a contrary direction, which motivated the double fuzzy implications-based restriction inference algorithm [19]. The variable differently implicational algorithm was further researched focusing on the FMT problem, in which the differently implicational principle for FMT was improved [20]. The continuous and uniformly continuous properties of the entropy-based differently implicational algorithm were demonstrated for the Tchebyshev and Hamming metrics [21]. The robustness becomes a hot research point to analyze the triple I method [22-23]. In conclusion, the triple I method shows several good properties, which includes strict logical basis, continuity, reversibility, robustness, and so on.

Regarding this topic, Pei proposed the triple I* method of FMT [24] from the perspective of another kind of reversibility, which focused on

$$(A(x) \rightarrow B(y)) \rightarrow (B^*(y) \rightarrow A^*(x)). \quad (5)$$

From a deeper viewpoint, the first and third fuzzy implications in (3) correspond to the implication connective in a logic system; and the second fuzzy implication in (3) reflects the “if-then” relation of fuzzy inference model. Based upon this idea, we extend (3) as follows:

$$(A(x) \rightarrow_1 B(y)) \rightarrow_2 (A^*(x) \rightarrow_1 B^*(y)), \quad (6)$$

where $\rightarrow_1, \rightarrow_2$ are two fuzzy implications [25-26]. The method derived from (6) is called the symmetric implicational method.

In this paper, we think about all of these formulas including (4), (5) and (6), then a new fuzzy inference method called the α -symmetric I* restriction method is proposed, which focuses on ($\alpha \in (0,1]$)

$$(A(x) \rightarrow_1 B(y)) \rightarrow_2 (B^*(y) \rightarrow_1 A^*(x)) < \alpha. \quad (7)$$

Moreover, to carefully control the reasoning process, we generalize it to

$$(A(x) \rightarrow_1 B(y)) \rightarrow_2 (B^*(y) \rightarrow_1 A^*(x)) < \alpha(x, y) \quad (8)$$

where $\alpha(x, y) \in (0,1]$ is a function with regard to x, y . The latter is called the $\alpha(x, y)$ -symmetric I* restriction method.

The aim of this study is to research the α -symmetric I* restriction method and the $\alpha(x, y)$ -symmetric I* restriction method.

2. Preliminaries

Definition 2.1. ([27]) If $\otimes : [0,1]^2 \rightarrow [0,1]$ satisfies the following conditions:

$$(i) a \otimes b = b \otimes a,$$

$$(ii) a \otimes (b \otimes c) = (a \otimes b) \otimes c,$$

$$(iii) 1 \otimes a = a,$$

$$(iv) a \leq b \text{ implies } a \otimes c \leq b \otimes c,$$

then \otimes is called a triangular norm (t-norm, for short) on $[0,1]$. If \otimes also satisfies $a \otimes \bigvee \{x_i \mid i \in P\} = \bigvee \{a \otimes x_i \mid i \in P\}$ ($a, x_i \in [0,1]$ and $P \neq \emptyset$), \otimes is said to be a left continuous t-norm.

Definition 2.2. ([27]) If $\oplus : [0,1]^2 \rightarrow [0,1]$ satisfies the following conditions:

$$(i) a \oplus b = b \oplus a,$$

$$(ii) a \oplus (b \oplus c) = (a \oplus b) \oplus c,$$

$$(iii) 0 \oplus a = a,$$

$$(iv) a \leq b \text{ implies } a \oplus c \leq b \oplus c,$$

then \oplus is called a triangular conorm (t-conorm, for short) on $[0,1]$.

Definition 2.3. ([28]) If $\rightarrow : [0,1]^2 \rightarrow [0,1]$ satisfies

$$0 \rightarrow 0 = 0 \rightarrow 1 = 1 \rightarrow 1 = 1, \quad 1 \rightarrow 0 = 0,$$

then \rightarrow is called a fuzzy implication on $[0,1]$. $a \rightarrow b$ can also be written as $I(a, b)$ ($a, b \in [0,1]$).

Definition 2.4. ([29]) Suppose that Z is any non-empty set, a mapping $C : Z \rightarrow [0,1]$ is defined as a fuzzy set on Z .

Definition 2.5. ([30]) If \otimes and \rightarrow are two mappings $[0,1]^2 \rightarrow [0,1]$, then (\otimes, \rightarrow) is called a residual pair, or \otimes, \rightarrow are residual to each other, if the following residual condition holds ($a, b, c \in [0,1]$):

$$a \otimes b \leq c \text{ if and only if } b \leq a \rightarrow c. \quad (9)$$

Lemma 2.1. ([14]) Let \otimes be a left continuous t-norm and ($a, b \in [0,1]$)

$$a \rightarrow b = \sup \{y \in [0,1] \mid a \otimes y \leq b\}, \quad (10)$$

then (\otimes, \rightarrow) is a residual pair, and \rightarrow satisfies

$$(C1) a \rightarrow b \text{ is increasing in the second variable,}$$

$$(C2) a \rightarrow b \text{ is right-continuous w.r.t. } b,$$

$$(C3) a \rightarrow b \text{ is decreasing in the first variable,}$$

$$(C4) a \leq b \text{ if and only if } a \rightarrow b = 1,$$

$$(C5) 1 \rightarrow a = a,$$

$$(C6) a \leq b \rightarrow c \text{ if and only if } b \leq a \rightarrow c,$$

$$(C7) a \rightarrow (b \rightarrow c) = b \rightarrow (a \rightarrow c),$$

$$(C8) \inf \{a \rightarrow x_i \mid i \in P\} = a \rightarrow \inf \{x_i \mid i \in P\},$$

$$(C9) \inf \{x_i \rightarrow b \mid i \in P\} = \sup \{x_i \mid i \in P\} \rightarrow b,$$

$$(C10) 0 \rightarrow a = 1,$$

$$(C11) a \rightarrow 1 = 1,$$

in which $a, b, c, x_i \in [0,1]$, P is not empty.

Definition 2.6. ([28]) Let \otimes be a left continuous t-norm and \rightarrow is obtained from (11), then \rightarrow is said to be an R-implication.

Definition 2.7. ([27]) A fuzzy negation is a decreasing function $N : [0,1] \rightarrow [0,1]$ which satisfies $N(0) = 1$, $N(1) = 0$. Furthermore, a fuzzy negation N is said to be

$$(i) \text{ strict if } N \text{ is continuous and strictly decreasing;}$$

$$(ii) \text{ strong if } N \text{ is an involution (i.e., } N(N(x)) = x \text{ for any } x \in [0,1].$$

$x \in [0, 1]$).

Definition 2.8. ([31]) A mapping $\rightarrow : [0, 1]^2 \rightarrow [0, 1]$ is said to be an (S, N)-implication if there exist a t-conorm \oplus and a fuzzy negation N such that $(a, b \in [0, 1])$

$$a \rightarrow b = N(a) \oplus b. \quad (11)$$

If N is a strong negation, then \rightarrow is said to be a strong implication (S-implication).

Lemma 2.2. ([32]) Let \rightarrow be an (S, N)-implication, then \rightarrow satisfies (C1), (C3), (C5), (C7), (C10), (C11).

Proposition 2.1. ([25]) Suppose that \rightarrow is a fuzzy implication satisfying (C1), (C2) and (C11), then the mapping $\otimes : [0, 1]^2 \rightarrow [0, 1]$ expressed by $(a, b \in [0, 1])$

$$a \otimes b = \inf\{x \in [0, 1] \mid b \leq a \rightarrow x\}, \quad (12)$$

is residual to \rightarrow .

Definition 2.9. ([29]) Suppose that Z is any non-empty set and that $F(Z)$ is the set of all fuzzy subsets on Z , the partial order relation \leq_F is defined as follows:

$$A \leq_F B \Leftrightarrow A(z_0) \leq B(z_0) \quad (A, B \in F(Z)).$$

Lemma 2.3. ([29]) $\langle F(Z), \leq_F \rangle$ is a complete lattice.

Example 2.1. Here are some familiar fuzzy implications, which include Lukasiewicz implication, Gödel implication, Goguen implication, Fodor implication and Kleene-Dienes implication $(a, b \in [0, 1], x' = 1 - x)$.

$$I_L(a, b) = \begin{cases} 1, & a \leq b \\ a' + b, & a > b \end{cases} \quad (\text{Lukasiewicz implication});$$

$$I_G(a, b) = \begin{cases} 1, & a \leq b \\ b, & a > b \end{cases} \quad (\text{Gödel implication});$$

$$I_{Go}(a, b) = \begin{cases} 1, & a = 0 \\ (b/a) \wedge 1, & a > 0 \end{cases} \quad (\text{Goguen implication});$$

$$I_{FD}(a, b) = \begin{cases} 1, & a \leq b \\ a' \vee b, & a > b \end{cases} \quad (\text{Fodor implication}).$$

$$I_{KD}(a, b) = a' \vee b \quad (\text{Kleene-Dienes implication}).$$

Their residual t-norm are respectively as follows:

$$a \otimes_L b = \begin{cases} a + b - 1, & a + b > 1 \\ 0, & a + b \leq 1 \end{cases},$$

$$a \otimes_G b = a \wedge b, \quad a \otimes_{Go} b = a \times b,$$

$$a \otimes_{FD} b = \begin{cases} a \wedge b, & a + b > 1 \\ 0, & a + b \leq 1 \end{cases}.$$

$$a \otimes_{KD} b = \begin{cases} b, & a + b > 1 \\ 0, & a + b \leq 1 \end{cases}.$$

Here I_{KD}, I_{FD}, I_{LK} Their residual t-norm are respectively as follows: are (S,N)-implications and $I_G, I_{Go}, I_{FD}, I_{LK}$ are R-implications.

3. The α -Symmetric I* Restriction Method

3.1. Basic Structure

Here we provide the basic structure of the α -symmetric I* restriction method.

Aiming at the FMT problem, from the viewpoint of the α -symmetric I* restriction method, we can achieve the following principle:

α -symmetric I* restriction principle: The conclusion A^* of FMT is the largest fuzzy set in $F(X)$ such that (7) holds for any $x \in X, y \in Y$.

Definition 3.1. Let $A \in F(X)$, $B, B^* \in F(Y)$, if A^* (in $F(X)$) lets (7) hold for any $x \in X, y \in Y$, then A^* is said to be an α -symmetric I* restriction solution.

Definition 3.2. Suppose that $A \in F(X)$, $B, B^* \in F(Y)$, and that nonempty set \mathbb{E} is the set of all α -symmetric I* restriction solutions, and finally that C^* (in $F(X)$) is the supremum of \mathbb{E} , then C^* is called an α -SupT-symmetric I* restriction solution.

Proposition 3.1 Let $\rightarrow_1, \rightarrow_2$ satisfies (C1). If C_1 is an α -symmetric I* restriction solution and

$$C_2 \leq_F C_1 \quad (C_1, C_2 \in F(X)),$$

then C_2 is also an α -symmetric I* restriction solution.

Proof. Since C_1 is an α -symmetric I* restriction solution, we have $(x \in X, y \in Y)$

$$(A(x) \rightarrow_1 B(y)) \rightarrow_2 (B^*(y) \rightarrow_1 C_1(x)) < \alpha.$$

Because $C_2 \leq_F C_1$ and $\rightarrow_1, \rightarrow_2$ satisfies (C1), one has $(x \in X, y \in Y)$

$$B^*(y) \rightarrow_1 C_2(x) \leq B^*(y) \rightarrow_1 C_1(x).$$

Thus we get

$$\alpha > (A(x) \rightarrow_1 B(y)) \rightarrow_2 (B^*(y) \rightarrow_1 C_1(x))$$

$$\geq (A(x) \rightarrow_1 B(y)) \rightarrow_2 (B^*(y) \rightarrow_1 C_2(x)).$$

As a result, C_2 is also an α -symmetric I* restriction solution.

End of the proof.

3.2. Optimal Solutions

Theorem 3.1 Suppose that $\rightarrow_1, \rightarrow_2$ satisfy (C1). Then there exists C as an α -symmetric I* restriction solution if and only

if the following formula holds for any $x \in X, y \in Y$:

$$(A(x) \rightarrow_1 B(y)) \rightarrow_2 (B^*(y) \rightarrow_1 0) < \alpha. \quad (13)$$

Proof. On the one hand, if (13) holds, then we can let $C(x) \equiv 0$ ($x \in X$). Here C obviously satisfies (7). Consequently, C is an α -symmetric I* restriction solution.

On the other hand, if there exists C as an α -symmetric I* restriction solution, then C satisfies (7), that is ($x \in X, y \in Y$)

$$(A(x) \rightarrow_1 B(y)) \rightarrow_2 (B^*(y) \rightarrow_1 C(x)) < \alpha.$$

Since $\rightarrow_1, \rightarrow_2$ satisfies (C1), one has ($x \in X, y \in Y$)

$$B^*(y) \rightarrow_1 0 \leq B^*(y) \rightarrow_1 C(x).$$

Hence we have

$$\alpha > (A(x) \rightarrow_1 B(y)) \rightarrow_2 (B^*(y) \rightarrow_1 C(x))$$

$$\geq (A(x) \rightarrow_1 B(y)) \rightarrow_2 (B^*(y) \rightarrow_1 0).$$

That is, (13) holds.

End of the proof.

Remark 3.1. Assume that (13) holds. As for an α -symmetric I* restriction solution C_1 (in $F(X)$), any fuzzy set C_2 (in $F(X)$) which is smaller than C_1 is also an α -symmetric I* restriction solution (from Proposition 3.1). This implies that there exist a lot of α -symmetric I* restriction solutions, which will include a particular solution, i.e.,

$$C_3(x) \equiv 1 \quad (x \in X).$$

Here C_3 is a strange solution, for which (7) always holds no matter what $A \rightarrow B$ and B^* are employed. As a result, we can find if the optimal α -symmetric I* restriction solution exists, then it should be the biggest one or the supremum of all solutions.

Theorem 3.2 If $\rightarrow_1, \rightarrow_2$ are R-implications, and (13) holds, and \otimes_1, \otimes_2 are the t-norms residual to $\rightarrow_1, \rightarrow_2$. Then the α -SupT-symmetric I* restriction solution is as follows ($x \in X$):

$$A^*(x) = \inf_{y \in Y} \{B^*(y) \otimes_1 ((A(x) \rightarrow_1 B(y)) \otimes_2 \alpha)\}. \quad (14)$$

Proof. To begin with, we let

$$G_1 = \{x \in X \mid A^*(x) = 0\}, \quad G_2 = \{x \in X \mid A^*(x) > 0\}.$$

Suppose that $C \in F(X)$ such that

$$C(x) = 0 \quad \text{for } x \in G_1$$

and

$$C(x) < A^*(x) \quad \text{for } x \in G_2.$$

Then we prove that C is an α -symmetric I* restriction

solution, i.e., the following formula holds for any $x \in X, y \in Y$:

$$(A(x) \rightarrow_1 B(y)) \rightarrow_2 (B^*(y) \rightarrow_1 C(x)) < \alpha.$$

If $x \in G_1$, then considering (13) holds, we have $C(x) = 0$ satisfies (7).

If $x \in G_2$, then considering (14) holds and $C(x) < A^*(x)$, we get

$$C(x) < B^*(y) \otimes_1 ((A(x) \rightarrow_1 B(y)) \otimes_2 \alpha) \quad (15)$$

holds for any $y \in Y$. We prove it by contradiction. Suppose that (7) does not hold. Then there exists $x_0 \in X$ and $y_0 \in Y$ (obviously $x_0 \in G_2$) such that

$$(A(x_0) \rightarrow_1 B(y_0)) \rightarrow_2 (B^*(y_0) \rightarrow_1 C(x_0)) \geq \alpha$$

holds. From residual condition, one has

$$\alpha \leq (A(x_0) \rightarrow_1 B(y_0)) \rightarrow_2 (B^*(y_0) \rightarrow_1 C(x_0)),$$

$$(A(x_0) \rightarrow_1 B(y_0)) \otimes_2 \alpha \leq B^*(y_0) \rightarrow_1 C(x_0),$$

$$B^*(y_0) \otimes_1 [(A(x_0) \rightarrow_1 B(y_0)) \otimes_2 \alpha] \leq C(x_0).$$

This contradicts (15). Consequently (7) holds, and thus C is an α -symmetric I* restriction solution.

Furthermore, we prove that A^* determined by (14) is the supremum of all α -symmetric I* restriction solutions.

Assume that $D \in F(X)$ and that there exists $x_0 \in X$ such that

$$D(x_0) > A^*(x_0).$$

Next we verify that D is not an α -symmetric I* restriction solution. In fact, it follows from (14) that there exists $y_0 \in Y$ such that

$$D(x_0) > B^*(y_0) \otimes_1 ((A(x_0) \rightarrow_1 B(y_0)) \otimes_2 \alpha).$$

We get from residual condition that

$$B^*(y_0) \otimes_1 ((A(x_0) \rightarrow_1 B(y_0)) \otimes_2 \alpha) < D(x_0),$$

$$(A(x_0) \rightarrow_1 B(y_0)) \otimes_2 \alpha \leq B^*(y_0) \rightarrow_1 D(x_0),$$

$$\alpha \leq (A(x_0) \rightarrow_1 B(y_0)) \rightarrow_2 (B^*(y_0) \rightarrow_1 D(x_0)).$$

As a result, D is not an α -symmetric I* restriction solution.

To sum up A^* determined by (14) is the supremum of all α -symmetric I* restriction solutions, i.e., the α -SupT-symmetric I* restriction solution.

End of the proof.

Lemma 3.1. Suppose that \rightarrow is an (S, N)-implications satisfying (C2), then the mapping $\otimes: [0, 1]^2 \rightarrow [0, 1]$ expressed

by (12) is residual to \rightarrow .

Proof. Since \rightarrow is an (S, N)-implications satisfying (C2), we get From Lemma 2.2 that \rightarrow satisfies (C1), (C2) and (C11). Then it follows from Proposition 2.1 that the mapping \otimes expressed by (12) is residual to \rightarrow .

End of the proof.

Theorem 3.3 If $\rightarrow_1, \rightarrow_2$ are (S, N)-implications satisfying (C2), and (13) holds, and \otimes_1, \otimes_2 are the operators residual to $\rightarrow_1, \rightarrow_2$. Then the α -SupT-symmetric I* restriction solution is expressed as (14).

Proof. From Lemma 3.1, it is similar to Theorem 3.2 that we can get the conclusion.

End of the proof.

Proposition 3.2 If $\rightarrow_1, \rightarrow_2 \in \{I_L, I_G, I_{Go}, I_{FD}\}$, and (13) holds, and \otimes_1, \otimes_2 are the t-norms residual to $\rightarrow_1, \rightarrow_2$. Then the α -SupT-symmetric I* restriction solution is expressed as (14).

Proof. Since I_L, I_G, I_{Go}, I_{FD} are R-implications, then we get the conclusion from Theorem 3.2.

End of the proof.

Proposition 3.3 If $\rightarrow_1, \rightarrow_2 \in \{I_{KD}, I_L, I_{FD}\}$, and (13) holds, and \otimes_1, \otimes_2 are the t-norms residual to $\rightarrow_1, \rightarrow_2$. Then the α -SupT-symmetric I* restriction solution is expressed as (14).

Proof. Since I_{KD}, I_L, I_{FD} are (S, N)-implications satisfying (C2), then we get the conclusion from Theorem 3.3.

End of the proof.

Example 3.1 Let $\rightarrow_1, \rightarrow_2$ respectively take I_G, I_{Go} . Suppose that (13) holds. Then the α -SupT-symmetric I* restriction solution is as follows ($x \in X$):

$$A^*(x) = \inf_{y \in Y} \{B^*(y) \wedge (I_G(A(x), B(y)) \times \alpha)\}.$$

When $\rightarrow_1 = \rightarrow_2$, the α -symmetric I* restriction method degenerates into the α -triple I* restriction method. We can obtain the following definitions and corollaries.

Definition 3.3. Let $A \in F(X)$, $B, B^* \in F(Y)$, if A^* (in $F(X)$) lets

$$(A(x) \rightarrow B(y)) \rightarrow (B^*(y) \rightarrow A^*(x)) < \alpha \quad (16)$$

hold for any $x \in X, y \in Y$, then A^* is said to be an α -triple I* restriction solution.

Definition 3.4. Suppose that $A \in F(X)$, $B, B^* \in F(Y)$, and that nonempty set \mathbb{F} is the set of all α -triple I* restriction solutions, and finally that C^* (in $F(X)$) is the supremum of \mathbb{F} , then C^* is called an α -SupT-triple I* restriction solution.

Corollary 3.1 Let \rightarrow satisfies (C1). If C_1 is an α -triple I* restriction solution and

$$C_2 \leq_F C_1 \quad (C_1, C_2 \in F(X)),$$

then C_2 is also an α -triple I* restriction solution.

Corollary 3.2 Suppose that \rightarrow satisfies (C1). Then there exists C as an α -symmetric I* restriction solution if and only

if the following formula holds for any $x \in X, y \in Y$:

$$(A(x) \rightarrow B(y)) \rightarrow (B^*(y) \rightarrow 0) < \alpha. \quad (17)$$

Corollary 3.3 If \rightarrow is an R-implications, and (17) holds, and \otimes is the t-norm residual to \rightarrow . Then the α -SupT-triple I* restriction solution is as follows ($x \in X$):

$$A^*(x) = \inf_{y \in Y} \{B^*(y) \otimes ((A(x) \rightarrow B(y)) \otimes \alpha)\}. \quad (18)$$

Corollary 3.4 If \rightarrow is an (S, N)-implications satisfying (C2), and (17) holds, and \otimes is the operator residual to \rightarrow . Then the α -SupT-symmetric I* restriction solution is expressed as (18).

Corollary 3.5 If $\rightarrow_1, \rightarrow_2 \in \{I_L, I_G, I_{Go}, I_{FD}, I_{KD}\}$, and (17) holds, and \otimes is the operator residual to \rightarrow . Then the α -SupT-symmetric I* restriction solution is expressed as (18).

4. The $\alpha(x,y)$ -Symmetric I* Restriction Method

4.1. Basic Structure

Here we show the basic structure of the $\alpha(x,y)$ -symmetric I* restriction method.

Focusing on the FMT problem, from the viewpoint of the $\alpha(x,y)$ -symmetric I* restriction method, we can obtain the following principle:

$\alpha(x,y)$ -symmetric I* restriction principle: The conclusion A^* of FMT is the largest fuzzy set in $F(X)$ such that (8) holds for any $x \in X, y \in Y$.

Definition 4.1. Let $A \in F(X)$, $B, B^* \in F(Y)$, if A^* (in $F(X)$) lets (8) hold for any $x \in X, y \in Y$, then A^* is said to be an $\alpha(x,y)$ -symmetric I* restriction solution.

Definition 4.2. Suppose that $A \in F(X)$, $B, B^* \in F(Y)$, and that nonempty set \mathbb{G} is the set of all $\alpha(x,y)$ -symmetric I* restriction solutions, and finally that C^* (in $F(X)$) is the supremum of \mathbb{G} , then C^* is called an $\alpha(x,y)$ -SupT-symmetric I* restriction solution.

Proposition 4.1 Let $\rightarrow_1, \rightarrow_2$ satisfies (C1). If C_1 is an $\alpha(x,y)$ -symmetric I* restriction solution and

$$C_2 \leq_F C_1 \quad (C_1, C_2 \in F(X)),$$

then C_2 is also an $\alpha(x,y)$ -symmetric I* restriction solution.

Proof. Because C_1 is an $\alpha(x,y)$ -symmetric I* restriction solution, we get ($x \in X, y \in Y$)

$$(A(x) \rightarrow_1 B(y)) \rightarrow_2 (B^*(y) \rightarrow_1 C_1(x)) < \alpha(x, y).$$

Since $C_2 \leq_F C_1$ and $\rightarrow_1, \rightarrow_2$ satisfies (C1), one has ($x \in X, y \in Y$)

$$B^*(y) \rightarrow_1 C_2(x) \leq B^*(y) \rightarrow_1 C_1(x).$$

Thus we obtain

$$\begin{aligned} \alpha(x, y) &> (A(x) \rightarrow_1 B(y)) \rightarrow_2 (B^*(y) \rightarrow_1 C_1(x)) \\ &\geq (A(x) \rightarrow_1 B(y)) \rightarrow_2 (B^*(y) \rightarrow_1 C_2(x)). \end{aligned}$$

Consequently, C_2 is also an $\alpha(x, y)$ -symmetric I* restriction solution.

End of the proof.

4.2. Optimal Solutions

Theorem 4.1 Suppose that $\rightarrow_1, \rightarrow_2$ satisfy (C1). Then there exists C as an $\alpha(x, y)$ -symmetric I* restriction solution if and only if the following formula holds for any $x \in X, y \in Y$:

$$(A(x) \rightarrow_1 B(y)) \rightarrow_2 (B^*(y) \rightarrow_1 0) < \alpha(x, y). \quad (19)$$

Proof. To begin with, if (19) holds, then we can let $C(x) \equiv 0$ ($x \in X$). Here C obviously satisfies (8). As a result, C is an $\alpha(x, y)$ -symmetric I* restriction solution.

What is more, if there exists C as an $\alpha(x, y)$ -symmetric I* restriction solution, then C satisfies (8), that is ($x \in X, y \in Y$)

$$(A(x) \rightarrow_1 B(y)) \rightarrow_2 (B^*(y) \rightarrow_1 C(x)) < \alpha(x, y).$$

Note that $\rightarrow_1, \rightarrow_2$ satisfies (C1), hence it follows that ($x \in X, y \in Y$)

$$B^*(y) \rightarrow_1 0 \leq B^*(y) \rightarrow_1 C(x).$$

Finally we obtain

$$\begin{aligned} \alpha(x, y) &> (A(x) \rightarrow_1 B(y)) \rightarrow_2 (B^*(y) \rightarrow_1 C(x)) \\ &\geq (A(x) \rightarrow_1 B(y)) \rightarrow_2 (B^*(y) \rightarrow_1 0). \end{aligned}$$

That is, (19) holds.

End of the proof.

Remark 4.1. Suppose that (19) holds. Considering an $\alpha(x, y)$ -symmetric I* restriction solution C_1 (in $F(X)$), any fuzzy set C_2 (in $F(X)$) which is smaller than C_1 is also an α -symmetric I* restriction solution (from Proposition 4.1). This means that there are a lot of $\alpha(x, y)$ -symmetric I* restriction solutions, which will incorporate a special solution, i.e.,

$$C_3(x) \equiv 1 \quad (x \in X).$$

Here C_3 is a particular solution, for which (8) always holds no matter what $A \rightarrow B$ and B^* are taken. As a result, we can know that if the optimal $\alpha(x, y)$ -symmetric I* restriction solution exists, then it should be the largest one or the supremum of all $\alpha(x, y)$ -symmetric I* restriction solutions.

Theorem 4.2 If $\rightarrow_1, \rightarrow_2$ are R-implications, and (19) holds, and \otimes_1, \otimes_2 are the t-norms residual to $\rightarrow_1, \rightarrow_2$. Then the $\alpha(x, y)$ -SupT-symmetric I* restriction solution is as follows

($x \in X$):

$$A^*(x) = \inf_{y \in Y} \{B^*(y) \otimes_1 ((A(x) \rightarrow_1 B(y)) \otimes_2 \alpha(x, y))\}. \quad (20)$$

Proof. First of all, we denote

$$H_1 = \{x \in X \mid A^*(x) = 0\}, \quad H_2 = \{x \in X \mid A^*(x) > 0\}.$$

Suppose that $C \in F(X)$ such that

$$C(x) = 0 \quad \text{for } x \in H_1$$

and

$$C(x) < A^*(x) \quad \text{for } x \in H_2.$$

Then we verify that C is an $\alpha(x, y)$ -symmetric I* restriction solution, i.e., the following formula holds for any $x \in X, y \in Y$:

$$(A(x) \rightarrow_1 B(y)) \rightarrow_2 (B^*(y) \rightarrow_1 C(x)) < \alpha(x, y).$$

If $x \in H_1$, then noting (19) holds, we have $C(x) = 0$ satisfies (8).

If $x \in H_2$, then noting (20) holds and $C(x) < A^*(x)$, we have

$$C(x) < B^*(y) \otimes_1 ((A(x) \rightarrow_1 B(y)) \otimes_2 \alpha(x, y)) \quad (21)$$

holds for any $y \in Y$. We prove it by contradiction. Suppose on the contrary that (8) does not hold. Then there exists $x_0 \in X$ and $y_0 \in Y$ (obviously $x_0 \in H_2$) such that

$$(A(x_0) \rightarrow_1 B(y_0)) \rightarrow_2 (B^*(y_0) \rightarrow_1 C(x_0)) \geq \alpha(x_0, y_0)$$

holds. From residual condition, one has

$$\alpha(x_0, y_0) \leq (A(x_0) \rightarrow_1 B(y_0)) \rightarrow_2 (B^*(y_0) \rightarrow_1 C(x_0)),$$

$$(A(x_0) \rightarrow_1 B(y_0)) \otimes_2 \alpha(x_0, y_0) \leq B^*(y_0) \rightarrow_1 C(x_0),$$

$$B^*(y_0) \otimes_1 [(A(x_0) \rightarrow_1 B(y_0)) \otimes_2 \alpha(x_0, y_0)] \leq C(x_0).$$

This contradicts (21). Consequently (8) holds, and thus C is an $\alpha(x, y)$ -symmetric I* restriction solution.

What is more, we show that A^* determined by (20) is the supremum of all $\alpha(x, y)$ -symmetric I* restriction solutions.

Suppose that $D \in F(X)$ and that there is $x_0 \in X$ such that

$$D(x_0) > A^*(x_0).$$

Next we verifies that D is not an $\alpha(x, y)$ -symmetric I* restriction solution. In fact, it follows from (20) that there exists $y_0 \in Y$ such that

$$D(x_0) > B^*(y_0) \otimes_1 ((A(x_0) \rightarrow_1 B(y_0)) \otimes_2 \alpha(x_0, y_0)).$$

It follows from residual condition that

$$B^*(y_0) \otimes_1 ((A(x_0) \rightarrow_1 B(y_0)) \otimes_2 \alpha(x_0, y_0)) < D(x_0),$$

$$(A(x_0) \rightarrow_1 B(y_0)) \otimes_2 \alpha(x_0, y_0) \leq B^*(y_0) \rightarrow_1 D(x_0),$$

$$\alpha(x_0, y_0) \leq (A(x_0) \rightarrow_1 B(y_0)) \rightarrow_2 (B^*(y_0) \rightarrow_1 D(x_0)).$$

Consequently, D is not an $\alpha(x, y)$ -symmetric I* restriction solution.

Summarizing above, A^* determined by (20) is the supremum of all $\alpha(x, y)$ -symmetric I* restriction solutions, i.e., the $\alpha(x, y)$ -SupT-symmetric I* restriction solution.

End of the proof.

Theorem 4.3 If $\rightarrow_1, \rightarrow_2$ are (S, N)-implications satisfying (C2), and (19) holds, and \otimes_1, \otimes_2 are the operators residual to $\rightarrow_1, \rightarrow_2$. Then the $\alpha(x, y)$ -SupT-symmetric I* restriction solution is expressed as (20).

Proof. From Lemma 3.1, it is similar to Theorem 4.2 that we can obtain the conclusion.

End of the proof.

Proposition 4.2 If $\rightarrow_1, \rightarrow_2 \in \{I_L, I_G, I_{Go}, I_{FD}\}$, and (19) holds, and \otimes_1, \otimes_2 are the t-norms residual to $\rightarrow_1, \rightarrow_2$. Then the $\alpha(x, y)$ -SupT-symmetric I* restriction solution is expressed as (20).

Proof. Because I_L, I_G, I_{Go}, I_{FD} are R-implications, then one has the conclusion from Theorem 4.2.

End of the proof.

Proposition 4.3 If $\rightarrow_1, \rightarrow_2 \in \{I_{KD}, I_L, I_{FD}\}$, and (19) holds, and \otimes_1, \otimes_2 are the t-norms residual to $\rightarrow_1, \rightarrow_2$. Then the $\alpha(x, y)$ -SupT-symmetric I* restriction solution is expressed as (20).

Proof. Because I_{KD}, I_L, I_{FD} are (S, N)-implications satisfying (C2), then we get the conclusion from Theorem 4.3.

End of the proof.

Example 4.1 Let $\rightarrow_1, \rightarrow_2$ respectively take I_{Go}, I_G . Suppose that (19) holds. Then the $\alpha(x, y)$ -SupT-symmetric I* restriction solution is as follows ($x \in X$):

$$A^*(x) = \inf_{y \in Y} \{B^*(y) \times (I_G(A(x), B(y)) \wedge \alpha(x, y))\}.$$

When $\rightarrow_1 = \rightarrow_2$, the $\alpha(x, y)$ -symmetric I* restriction method degenerates into the $\alpha(x, y)$ -triple I* restriction method. We can obtain the following definitions and corollaries.

Definition 4.3. Let $A \in F(X)$, $B, B^* \in F(Y)$, if A^* (in $F(X)$) lets

$$(A(x) \rightarrow B(y)) \rightarrow (B^*(y) \rightarrow A^*(x)) < \alpha(x, y) \quad (22)$$

hold for any $x \in X, y \in Y$, then A^* is said to be an $\alpha(x, y)$ -triple I* restriction solution.

Definition 4.4. Suppose that $A \in F(X)$, $B, B^* \in F(Y)$, and that nonempty set \mathbb{H} is the set of all $\alpha(x, y)$ -triple I* restriction solutions, and finally that C^* (in $F(X)$) is the supremum of \mathbb{H} , then C^* is called an $\alpha(x, y)$ -SupT-triple I*

restriction solution.

Corollary 4.1 Let \rightarrow satisfies (C1). If C_1 is an $\alpha(x, y)$ -triple I* restriction solution and

$$C_2 \leq_F C_1 \quad (C_1, C_2 \in F(X)),$$

then C_2 is also an $\alpha(x, y)$ -triple I* restriction solution.

Corollary 4.2 Suppose that \rightarrow satisfies (C1). Then there exists C as an $\alpha(x, y)$ -symmetric I* restriction solution if and only if the following formula holds for any $x \in X, y \in Y$:

$$(A(x) \rightarrow B(y)) \rightarrow (B^*(y) \rightarrow 0) < \alpha(x, y). \quad (23)$$

Corollary 4.3 If \rightarrow is an R-implications, and (23) holds, and \otimes is the t-norm residual to \rightarrow . Then the $\alpha(x, y)$ -SupT-triple I* restriction solution is as follows ($x \in X$):

$$A^*(x) = \inf_{y \in Y} \{B^*(y) \otimes ((A(x) \rightarrow B(y)) \otimes \alpha(x, y))\}. \quad (24)$$

Corollary 4.4 If \rightarrow is an (S, N)-implications satisfying (C2), and (23) holds, and \otimes is the operator residual to \rightarrow . Then the $\alpha(x, y)$ -SupT-symmetric I* restriction solution is expressed as (24).

Corollary 4.5 If $\rightarrow \in \{I_L, I_G, I_{Go}, I_{FD}, I_{KD}\}$, and (23) holds, and \otimes is the operator residual to \rightarrow . Then the $\alpha(x, y)$ -SupT-symmetric I* restriction solution is expressed as (24).

5. Conclusions

We consider the idea of symmetric implicational reasoning, triple I* method and restriction theory, then we put forward the α -symmetric I* restriction method and the $\alpha(x, y)$ -symmetric I* restriction method.

The main results and conclusions are as follows:

- First of all, to find the optimal solutions, the α -symmetric I* restriction principle and the $\alpha(x, y)$ -symmetric I* restriction principle are revealed.
- Moreover, solution is the key for a fuzzy inference method. The unified solutions of the two methods are achieved, especially for R-implications and (S, N)-implications.
- Finally, from different viewpoint, when $\rightarrow_1 = \rightarrow_2$, the two methods degenerate into the α -triple I* restriction method and $\alpha(x, y)$ -triple I* restriction method. Then corresponding conclusions are given.

In the next step, we shall extend the proposed methods under the environment of granular computing [33-34].

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